

COHOMOLOGY GROUPS OF NON-UNIFORM RANDOM SIMPLICIAL COMPLEXES

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ABSTRACT. We consider a model of a random simplicial complex generated by taking the downward-closure of a non-uniform binomial random hypergraph, in which each set of $k + 1$ vertices forms an edge with some probability p_k independently, where p_k depends on k and on the number of vertices n . We consider a notion of connectedness on this model according to the vanishing of cohomology groups over an arbitrary abelian group R . We prove that this notion of connectedness displays a phase transition and determine the threshold. We also prove a hitting time result for a natural process interpretation, in which simplices and their downward-closure are added one by one.

1. INTRODUCTION

1.1. Motivation

One of the first and most famous results in the theory of random graphs, due to Erdős and Rényi [5], states that the uniform random graph $G(n, m)$ displays a phase transition threshold for the property of being connected at about $m = \frac{1}{2}n \log n$. Almost equivalently, in modern terminology, the binomial random graph $G(n, p)$ becomes connected around $p = \frac{\log n}{n}$. The result was subsequently strengthened by Bollobás and Thomason [1] to a *hitting time* result – the random graph process, in which edges are added to an empty graph one by one in a uniformly random order, is very likely to become connected at exactly the moment at which the last isolated vertex disappears (i.e. acquires an edge), which occurs around time $p = \frac{\log n}{n}$.

There have been many generalisations of these results to higher-dimensional analogues of graphs, including hypergraphs (e.g. [3, 7]) and simplicial complexes (e.g. [2, 7, 8, 9])

In this paper we prove a generalisation for a model of simplicial complexes arising from non-uniform binomial random hypergraphs. The notion of connectedness that we study regards the vanishing of cohomology groups over an abelian group.

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This extends recent results of [2] in two important ways: firstly in that the random hypergraph used to generate the random complex is *non-uniform*, having different probabilities for each edge size; and secondly in that we consider cohomology groups over *any* abelian group R , rather than just over \mathbb{F}_2 . It is also closely related to results of Linial and Meshulam [8], Meshulam and Wallach [9], and Kahle and Pittel [7], although the model of random simplicial complexes that we consider is different and leads to significantly more complex behaviour.

1.2. Model

Throughout the paper let $d \geq 2$ be a fixed integer and let R be an abelian group with at least two elements. We use additive notation for the group operation of R and denote its neutral element by 0_R . Given an integer k , we write $[k] := \{1, \dots, k\}$.

We define a model of a random d -complex generated from a non-uniform random hypergraph, in which sets of vertices have different probabilities of forming an edge depending on their size.

Definition 1.1. For each $k \in [d]$, let $p_k = p_k(n) \in [0, 1] \subset \mathbb{R}$ be given and write $\mathbf{p} := (p_1, \dots, p_d)$. Denote by $G_{\mathbf{p}} = G(n, \mathbf{p})$ the (non-uniform) binomial random hypergraph on vertex set $[n]$ in which, for all $k \in [d]$, each element of $\binom{[n]}{k+1}$ forms a hyperedge with probability p_k independently. By $\mathcal{G}_{\mathbf{p}} = \mathcal{G}(n, \mathbf{p})$, we denote the random d -dimensional simplicial complex on $[n]$ such that

- the 0-simplices of $\mathcal{G}_{\mathbf{p}}$ are the singletons of $[n]$ and
- for each $i \in [d]$, the i -simplices are precisely the $(i+1)$ -sets which are contained in hyperedges of $G_{\mathbf{p}}$.

In other words, $\mathcal{G}_{\mathbf{p}}$ is the downward-closure of the set of hyperedges of $G_{\mathbf{p}}$, together with all singletons of $[n]$ (if those are not already in the downward-closure).

Recently, a similar but slightly more general model was independently introduced in [6]. This model has a similar flavour to the multi-parameter model introduced by Costa and Farber [4].

Denote by $H^i(\mathcal{G}; R)$ the i -th cohomology group of a simplicial complex \mathcal{G} with coefficients in R . It is well-known that $H^0(\mathcal{G}; R) = R$ if and only if \mathcal{G} is connected in the topological sense (see e.g. [10, Theorem 42.1]), which we call *topologically connected* in order to distinguish it from other notions of connectedness. For any positive integer i , the *vanishing* of $H^i(\mathcal{G}; R)$ can be viewed as a “higher-order connectedness” of \mathcal{G} .

Definition 1.2. Given a positive integer j , a simplicial complex \mathcal{G} is called *R -cohomologically j -connected* (*j -cohom-connected* for short) if

- (i) $H^0(\mathcal{G}; R) = R$;
- (ii) $H^i(\mathcal{G}; R) = 0$ for all $i \in [j]$.

1.3. Main results

The main aim of this paper is to provide an analogue of the graph results of Erdős and Rényi [5] and of Bollobás and Thomason [1] for the $\mathcal{G}_{\mathbf{p}}$ model of random

simplicial complexes. Given an appropriate direction $\bar{\mathbf{p}}$, we consider the random simplicial complex process $(\mathcal{G}_\tau) = (\mathcal{G}_{\tau\bar{\mathbf{p}}})_{\tau \in \mathbb{R}_{\geq 0}}$ (defined more formally later) and describe for which values of τ the complex is j -cohom-connected and for which it is not, relating this threshold to the disappearance of the last minimal obstruction.

In order to define the minimal obstructions \hat{M}_j^k , we introduce the following necessary concepts.

Definition 1.3. Let $j \in [d-1]$ and let k be an integer with $j+1 \leq k \leq d$. Given a k -simplex K in a d -dimensional simplicial complex \mathcal{G} , we say that a collection $\mathcal{F} = \{P_0, \dots, P_{k-j}\}$ of j -simplices forms a j -flower in K (see Figure 1) if

- (F1) $K = \bigcup_{i=0}^{k-j} P_i$;
- (F2) $C := \bigcap_{i=0}^{k-j} P_i$ satisfies $|C| = j$.

We call the j -simplices P_i the *petals* and the set C the *centre* of the j -flower \mathcal{F} .

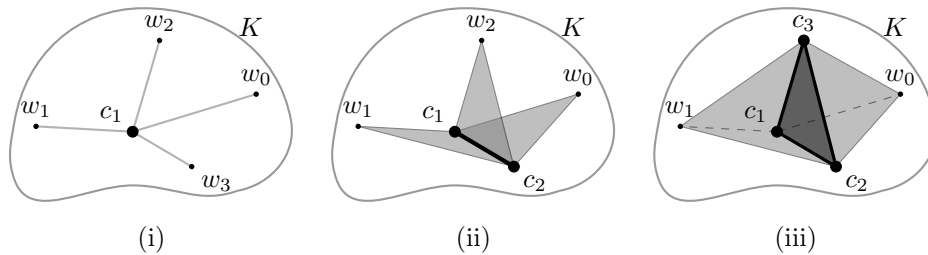


Figure 1. Examples of j -flowers in a k -simplex K , for $k = 4$ and $j = 1, 2, 3$.

- (i) The 1-flower in K with centre $C = \{c_1\}$ (bold black) and petals $P_i = C \cup \{w_i\}$, $i = 0, 1, 2, 3$ (grey).
- (ii) The 2-flower in K with centre $C = \{c_1, c_2\}$ (bold black) and petals $P_i = C \cup \{w_i\}$, $i = 0, 1, 2$ (grey).
- (iii) The 3-flower in K with centre $C = \{c_1, c_2, c_3\}$ (bold black) and petals $P_i = C \cup \{w_i\}$, $i = 0, 1$ (grey).

Observe that for each k -simplex K and each $(j - 1)$ -simplex $C \subseteq K$, there is a unique j -flower in K with centre C , namely

$$(1) \quad \mathcal{F}(K, C) := \{C \cup \{w\} \mid w \in K \setminus C\}.$$

Definition 1.4. For any $(j + 2)$ -set A in a complex \mathcal{G} , the collection of all $(j + 1)$ -sets of A is called a j -shell if each of them forms a j -simplex in \mathcal{G} .

If the collection of all $(j + 1)$ -subsets of a $(j + 2)$ -set A forms a j -shell, with a slight abuse of terminology we also refer to the set A itself as a j -shell.

Definition 1.5. For $j + 1 \leq k \leq d$, a 4-tuple (K, C, w, a) is said to form a *copy* of \hat{M}_j^k if

- (M1) K is a k -simplex in \mathcal{G} ;
- (M2) C is a $(j - 1)$ -simplex in K such that every simplex of \mathcal{G} that contains a petal of the flower $\mathcal{F} = \mathcal{F}(K, C)$ is contained in K ;

(M3) $w \in K \setminus C$ and $a \in [n] \setminus K$ are such that $C \cup \{w\} \cup \{a\}$ is a j -shell in \mathcal{G} .

We call the j -simplex $C \cup \{w\}$ the *base* and a the *apex vertex* of the j -shell $C \cup \{w\} \cup \{a\}$. Every other j -simplex in $C \cup \{w\} \cup \{a\}$ is called a *side* of the j -shell.

A copy of \hat{M}_j^k is an obstruction to the vanishing of $H^j(\mathcal{G}; R)$, since it is easy to define a *bad function*, i.e. a j -cocycle which is not a j -coboundary, by choosing any non-zero element $r \in R \setminus \{0_R\}$ and assigning the values $\pm r$ to the (ordered) petals in an appropriate way. Furthermore, it can be shown that a copy of \hat{M}_j^k is a *minimal* obstruction in terms of the size of the support of a bad function.

The random d -complex $\mathcal{G}_{\mathbf{p}}$ can be turned into a *process*, by assigning a *birth time* to each k -simplex. More precisely, for each $k \in [d]$ and each $(k + 1)$ -set $K \in \binom{[n]}{k+1}$ independently, sample a birth time uniformly at random from $[0, 1]$. Then $\mathcal{G}_{\mathbf{p}}$ is exactly the complex generated by the $(k + 1)$ -sets with birth times at most p_k , for all $k \in [d]$, by taking the downward-closure. If we fix a “direction” $\bar{\mathbf{p}} = (\bar{p}_1, \dots, \bar{p}_d)$ of non-negative real numbers with $\bar{p}_d \neq 0$, set

$$\mathbf{p} = \tau \bar{\mathbf{p}} := (\min\{\tau p_1, 1\}, \dots, \min\{\tau p_d, 1\}),$$

and gradually increase τ from 0 to $\tau_{\max} := 1/\bar{p}_d$, then $\mathcal{G}_{\mathbf{p}}$ becomes a process in which simplices (together with their downward-closure) arrive one by one. We will denote this process by $(\mathcal{G}_{\tau \bar{\mathbf{p}}})_{\tau}$, or sometimes just by (\mathcal{G}_{τ}) when the direction $\bar{\mathbf{p}}$ is clear from the context.

Given a direction $\bar{\mathbf{p}}$, and a $(k + 1)$ -set K with birth time t_K , the *scaled birth time* of K is

$$\tau_K := \frac{t_K}{\bar{p}_k}.$$

Thus τ_K is distributed uniformly in $[0, \bar{p}_k^{-1}]$, and $\mathcal{G}_{\tau} = \mathcal{G}_{\tau \bar{\mathbf{p}}}$ consists of all those simplices with *scaled* birth time at most τ , together with their downward-closure.

With some elementary arguments we can show that, when considering j -cohom-connectedness, it suffices to consider a direction $\bar{\mathbf{p}}$ with some specific properties, namely that for each $1 \leq k \leq d$ there are constants $\bar{\alpha}_k, \bar{\gamma}_k$ and a function $\bar{\beta}_k = \bar{\beta}_k(n)$ such that

$$(2) \quad \bar{p}_k = \frac{\bar{\alpha}_k \log n + \bar{\beta}_k}{n^{k-j+\bar{\gamma}_k}} (k - j)!,$$

and furthermore

- (P1) at least one of $\bar{\alpha}_k, \bar{\gamma}_k$ is zero and neither of them is negative;
- (P2) if $\bar{\alpha}_k = 0$, then $\bar{\beta}_k$ is either zero or it is positive and *subpolynomial* in the sense that for every constant $\varepsilon > 0$, we have $\bar{\beta}_k = o(n^\varepsilon)$, but $\bar{\beta}_k = \omega(n^{-\varepsilon})$;
- (P3) if $\bar{\gamma}_k = 0$, then $|\bar{\beta}_k| = o(\log n)$;
- (P4) there exists an index $j + 1 \leq k_0 \leq d$ with $\bar{\alpha}_{k_0} > 0$.

Additionally, for all indices k with $j \leq k \leq d$ and $\bar{p}_k \neq 0$, if we define the parameters

$$\begin{aligned}
 \bar{\lambda}_k &:= j + 1 - \bar{\gamma}_k - (k - j + 1) \sum_{i=j+1}^d \bar{\alpha}_i, \\
 \bar{\mu}_k &:= -(k - j + 1) \sum_{i=j+1}^d \frac{\bar{\beta}_i}{n^{\bar{\gamma}_i}} + \begin{cases} \log \log n & \text{if } \bar{\alpha}_k \neq 0, \\ \log(\bar{\beta}_k) & \text{if } \bar{\alpha}_k = 0, \end{cases} \\
 \bar{\nu}_k &:= \begin{cases} -\log((j + 1)!) & \text{if } k = j, \\ -\log(j!) - \log(k - j + 1) + \log(\bar{\alpha}_k) & \text{if } \bar{\alpha}_k \neq 0, \\ -\log(j!) - \log(k - j + 1) & \text{otherwise,} \end{cases}
 \end{aligned}
 \tag{3}$$

then we may assume that

- (C1) $\bar{\lambda}_k \log n + \bar{\mu}_k + \bar{\nu}_k \leq 0$, for all indices k with $j \leq k \leq d$ and $\bar{p}_k \neq 0$,
- (C2) $\bar{\lambda}_{\bar{k}} \log n + \bar{\mu}_{\bar{k}} + \bar{\nu}_{\bar{k}} = 0$, for some \bar{k} with $j \leq \bar{k} \leq d$.

The definitions of the parameters $\bar{\lambda}_k, \bar{\mu}_k, \bar{\nu}_k$ are motivated by Lemma 2.2. We call a direction $\bar{\mathbf{p}}$ satisfying (P1)–(P4) and (C1)–(C2) a j -critical direction.

Theorem 1.6. For $j \in [d - 1]$ and a j -critical direction $\bar{\mathbf{p}}$, let $\mathbf{p}^* = \tau^* \bar{\mathbf{p}}$, where

$$\tau^* := \sup\{\tau \in \mathbb{R}_{\geq 0} \mid \mathcal{G}_{\tau \bar{\mathbf{p}}} \text{ contains a copy of } \hat{M}_j^k \text{ for some } j \leq k \leq d\}.$$

Then for every function ω of n which tends to infinity as $n \rightarrow \infty$, the following statements hold with high probability.

- (i) $\tau^* = 1 + o\left(\frac{\omega}{\log n}\right)$.
- (ii) The random d -complex process $(\mathcal{G}_\tau) = (\mathcal{G}_{\tau \bar{\mathbf{p}}})_\tau$ is not R -cohomologically j -connected for all $\tau < \tau^*$, i.e.

$$H^0(\mathcal{G}_{\mathbf{p}}; R) \neq R \quad \text{or} \quad H^i(\mathcal{G}_{\mathbf{p}}; R) \neq 0 \text{ for some } i \in [j],$$

for $\mathbf{p} = \tau \bar{\mathbf{p}}$ and for $\tau \in [0, \tau^*)$.

- (iii) The random d -complex process $(\mathcal{G}_\tau) = (\mathcal{G}_{\tau \bar{\mathbf{p}}})_\tau$ is R -cohomologically j -connected for all $\tau \geq \tau^*$, i.e.

$$H^0(\mathcal{G}_{\mathbf{p}}; R) = R \quad \text{and} \quad H^i(\mathcal{G}_{\mathbf{p}}; R) = 0 \text{ for all } i \in [j],$$

for $\mathbf{p} = \tau \bar{\mathbf{p}}$ and for $\tau \in [\tau^*, \tau_{\max}]$.

Let us note that neither j -cohom-connectedness nor the presence of copies of \hat{M}_j^k are necessarily monotone properties, which makes the proofs significantly harder. Indeed, it is not immediately clear that j -cohom-connectedness should have a single threshold – in principle, $\mathcal{G}_\tau = \mathcal{G}_{\tau \bar{\mathbf{p}}}$ could switch between being j -cohom-connected or not several times. However, Theorem 1.6 implies that with high probability it does not.

2. MINIMAL OBSTRUCTIONS AND THE HITTING TIME

Let us now define a ‘reduced’ version of \hat{M}_j^k , denoted by M_j^k , by omitting the condition (M3) on the j -shell $C \cup \{w\} \cup \{a\}$ in Definition 1.5. It is easy to see that this shell is very likely to exist if τ is ‘large enough’ ($\tau \geq \varepsilon/(\log n)$ will do), which will be the case well before the critical range for the disappearance of M_j^k . Thus it is convenient to switch attention from \hat{M}_j^k to M_j^k .

Definition 2.1. Let k be an integer with $j + 1 \leq k \leq d$. A pair (K, C) is called a *copy of M_j^k* if it satisfies the following conditions.

- (M1) K is a k -simplex in \mathcal{G} ;
- (M2) C is a $(j - 1)$ -simplex in K such that every simplex of \mathcal{G} that contains a petal of the flower $\mathcal{F} = \mathcal{F}(K, C)$ is contained in K .

For $k = j$, an isolated j -simplex is called a copy of M_j^j and of \hat{M}_j^j .

Lemma 2.2. *Suppose $\tau = O(1)$ and $\bar{\mathbf{p}}$ is a direction vector satisfying (P1)–(P4). Then the number X_j^k of copies of M_j^k in \mathcal{G}_τ satisfies*

$$(4) \quad \log \mathbb{E}(X_j^k) = \lambda_k \log n + \mu_k + \nu_k + o(1)$$

for all $j \leq k \leq d$ with $p_k \neq 0$, where λ_k, μ_k, ν_k are defined analogously to (3) for $\mathbf{p} = \tau \bar{\mathbf{p}}$ instead of $\bar{\mathbf{p}}$.

Since $\bar{\mathbf{p}}$ in Theorem 1.6 was chosen to be j -critical, i.e. such that the right hand side of (4) is at most $o(1)$ for all k and for $\mathbf{p} = \bar{\mathbf{p}}$, this tells us that heuristically, the last minimal obstruction should disappear around time $\tau = 1$.

Indeed, let τ' denote the smallest scaled birth time $\tau \geq 1 - \frac{\log \log n}{10d \log n}$ such that \mathcal{G}_τ is M_j^k -free for all k . Extending Lemma 2.2 with a second moment argument shows that τ' is close to 1 with high probability. Furthermore, we can also show that new copies of M_j^k are unlikely to appear after time $1 - o(1)$. Together with the fact that the existence of M_j^k and \hat{M}_j^k are essentially equivalent events, we obtain the following corollary.

Corollary 2.3. *Let $\omega = \omega(n)$ be a function that tends to infinity as $n \rightarrow \infty$. Then with high probability*

$$1 - \frac{\omega}{\log n} < \tau^* = \tau' < 1 + \frac{\omega}{\log n}.$$

3. COVERING THE SUBCRITICAL CASE

To prove statement (ii) of Theorem 1.6, we first observe that initially the process is not topologically connected.

Lemma 3.1. *There exist positive constants $c^- = c^-(d)$ and $c^+ = c^+(d)$ such that*

- (i) *Whp $\mathcal{G}_\mathbf{p}$ is not topologically connected if $p_i \leq \frac{c^- \log n}{n^i}$ for all $i \in [d]$;*
- (ii) *Whp $\mathcal{G}_\mathbf{p}$ is topologically connected if $p_i \geq \frac{c^+ \log n}{n^i}$ for some $i \in [d]$.*

The proof is an elementary extension of the graph case: for (i) we use a second moment method to show that with high probability there are isolated vertices. Statement (ii) can be proved analogously to the graph case, or follows as a special case of far stronger results in [3] or [11].

Subsequently, we proceed by induction on j , and the induction hypothesis tells us that up to time $\tau = \frac{\Theta(1)}{n}$, whp the process is not even $(j - 1)$ -cohom-connected. (Lemma 3.1 effectively provides the base case $j = 0$ of the induction.) It remains to cover the range $\tau \in [\varepsilon/n, \tau^*)$ for some sufficiently small constant ε . We achieve this by splitting the range into three subintervals.

Lemma 3.2. *For every constant $\varepsilon > 0$ there exists a constant $\delta > 0$ such that the following is true. Let*

$$I_1 := \left[\frac{\varepsilon}{n}, \frac{\delta}{n} \right], \quad I_2 := \left[\frac{\delta}{n}, 1 - \frac{1}{(\log n)^{1/3}} \right], \quad I_3 := \left[1 - \frac{1}{(\log n)^{1/3}}, \tau^* \right].$$

Then with high probability there exist k_1, k_2, k_3 and for each $i = 1, 2, 3$ a copy of $\hat{M}_j^{k_i}$ that is present throughout the range I_i .

The proof strategy is to show, via a second moment argument, that there are many copies of \hat{M}_j^k for some k at times $\tau = \varepsilon/n$ and $\tau = 1 - \frac{1}{(\log n)^{1/3}}$, and therefore with high probability at least one of these survives until the end of I_1 or was born by the start of I_2 respectively. On the other hand, since $\tau^* = 1 + o(1)$ whp, the last copy of \hat{M}_j^k to disappear at time τ^* was already present at any time $\tau = 1 - o(1)$ whp, and therefore whp exists throughout the range I_3 .

4. SUPERCRITICAL CASE

To prove statement (iii) of Theorem 1.6, we need to show two things: firstly, that whp \mathcal{G}_τ is j -cohom-connected at time $\tau = \tau^*$, and secondly that whp it does not become disconnected later.

In order for \mathcal{G}_τ not to be j -cohom-connected, it would have to admit a j -cocycle f_τ which is not a coboundary: we call such a function f_τ a *bad function*, and denote its support by \mathcal{S}_τ . Indeed, we may assume that \mathcal{S}_τ is the smallest possible support among all such bad functions in \mathcal{G}_τ .

The next lemma immediately implies that for any single choice of $\tau \geq \tau^*$, whp \mathcal{G}_τ is j -cohom-connected.

Lemma 4.1. *Suppose that $\tau \geq \tau^*$.*

- (i) *Let $h \in \mathbb{R}$ be any large constant. Then with high probability $|\mathcal{S}_\tau| > h$, if it exists.*
- (ii) *There exists a constant $\bar{h} \in \mathbb{R}$ such that with high probability $|\mathcal{S}_\tau| < \bar{h}$, if it exists.*

Indeed, we even prove a little more than this, showing that if $\tau^* > \tau = 1 - o(1)$, then the only obstructions to j -cohom-connectedness are copies of \hat{M}_j^k .

The proof of this lemma exploits a very useful property of such a minimal support \mathcal{S}_τ which we call *traversability*, and which in particular implies that \mathcal{S}_τ

can be explored via a search process. This allows us to bound the number of possible configurations for \mathcal{S}_τ . In case (ii) we also use a result from [9] to give a lower bound on the number of simplices which must *not* be present in order for \mathcal{S}_τ to form the support of a j -cocycle.

Finally, to prove that whp the process (\mathcal{G}_τ) never becomes disconnected again at any time $\tau > \tau^*$, we first prove that in order for a bad function f_τ to appear with the birth of a simplex K , its support \mathcal{S}_τ would have to be K -localised in the sense that within \mathcal{G}_{τ^*} any simplex containing an element of \mathcal{S}_τ must lie entirely within K . We then show that with high probability in \mathcal{G}_{τ^*} there are not many candidate positions for where such a simplex K might be born, and the probability that any one of these is born before any of the simplices that would prevent \mathcal{S}_τ from being K -localised is very small, so whp no new bad function will appear.

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