

*XIV: Numerical methods for pricing  
American options - PSOR algorithm*

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# Numerical solution

- Recall:

$$\left( \frac{\partial u}{\partial \tau} - \frac{\sigma^2}{2} \frac{\partial^2 u}{\partial x^2} \right) (u(x, \tau) - g(x, \tau)) = 0,$$

$$\frac{\partial u}{\partial \tau} - \frac{\sigma^2}{2} \frac{\partial^2 u}{\partial x^2} \geq 0, \quad u(x, \tau) - g(x, \tau) \geq 0$$

- Discretization - in the same way as in the implicit scheme for European options:

$$\mathbf{A}u^{j+1} \geq u^j + b^j, \quad u^{j+1} \geq g^{j+1} \quad \text{for } j = 0, 1, \dots, m-1,$$
$$(\mathbf{A}u^{j+1} - u^j - b^j)_i (u^{j+1} - g^{j+1})_i = 0 \quad \forall i$$

# Numerical solution

- Matrix  $A$  and vector  $b$  remain the same:

$$\mathbf{A} = \begin{pmatrix} 1 + 2\gamma & -\gamma & 0 & \cdots & 0 \\ -\gamma & 1 + 2\gamma & -\gamma & & \vdots \\ 0 & \cdot & \cdot & \cdot & 0 \\ \vdots & & & -\gamma & 1 + 2\gamma & -\gamma \\ 0 & \cdots & 0 & -\gamma & 1 + 2\gamma \end{pmatrix},$$

$$b^j = (\gamma\phi^{j+1}, 0, \dots, 0, \gamma\psi^{j+1})^T$$

where  $\gamma = \frac{\sigma^2 k}{2h^2}$

# PSOR method

- On each time level we solve a problem of the form

$$\begin{aligned} \mathbf{A}u &\geq b, & u &\geq g, \\ (\mathbf{A}u - b)_i (u_i - g_i) &= 0 & \forall i. \end{aligned}$$

- Define the sequence

$$u^0 = 0, \quad u^{p+1} = \max(\mathbf{T}_\omega u^p + c_\omega, g) \quad \text{for } p = 1, 2, \dots,$$

where  $T_\omega, c_\omega$  come from the classical SOR method and maximum is taken componentwise

- Projected SOR  $\rightarrow$  known as PSOR method or PSOR algorithm

# PSOR method

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- Components of the approximate solution:

$$u_i^{p+1} = \max \left[ \frac{\omega}{A_{ii}} \left( b_i - \sum_{j < i} A_{ij} u_j^{p+1} - \sum_{j > i} A_{ij} u_j^p \right) + (1 - \omega) u_i^p, g_i \right]$$

# Convergence of the algorithm to the solution

- The sequence  $u^p$  converges to some limit  $u$  - proof uses Banach fixed point theorem [Ševčovič, Stehlíková, Mikula: **Analytical and numerical methods for pricing financial derivatives**, pp. 156-157]
- This limit is a solution:
  - $u_i^{p+1} \geq g_i \Rightarrow$  also the limit satisfies  $u_i \geq g_i$
  - $u_i^{p+1} \geq \frac{\omega}{A_{ii}} \left( b_i - \sum_{j<i} A_{ij} u_j^{p+1} - \sum_{j>i} A_{ij} u_j^p \right) + (1 - \omega) u_i^p \Rightarrow$  also the limit satisfies  $u_i \geq \frac{\omega}{A_{ii}} \left( b_i - \sum_{j<i} A_{ij} u_j - \sum_{j>i} A_{ij} u_j^p \right) + (1 - \omega) u_i$   
we use that  $A_{ii} > 0, \omega > 0 \rightarrow$  we obtain  $(Au)_i \geq b_i$
  - if  $u_i > g_i$ , then starting with some index  $p_0$  we have  $u_i^p > g_i$ ; for these indices:

## Convergence of the algorithm to the solution

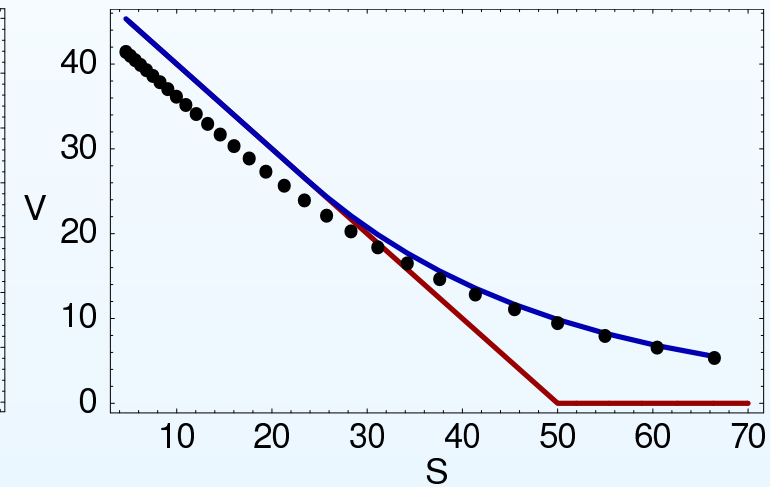
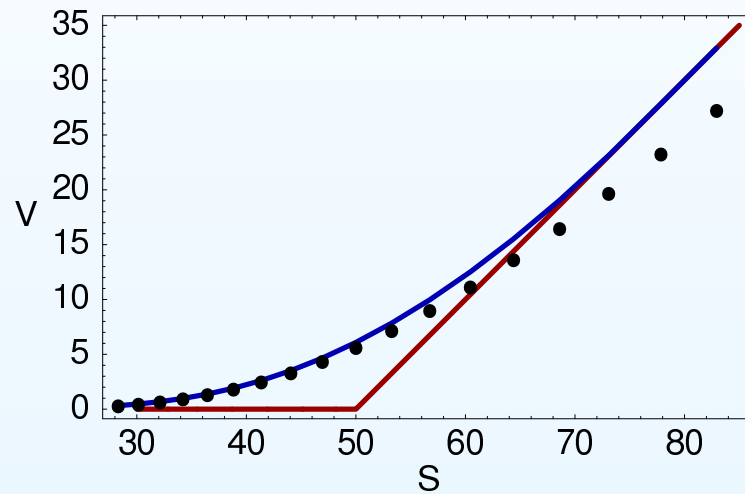
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$$u_i^{p+1} = \frac{\omega}{A_{ii}} \left( b_i - \sum_{j < i} A_{ij} u_j^{p+1} - \sum_{j > i} A_{ij} u_j^p \right) + (1 - \omega) u_i^p,$$

taking limit as  $p \rightarrow \infty$  we get  $(Au)_i = b_i \Rightarrow$  condition  $(Au - b)_i (u_i - g_i) = 0$  is satisfied

# Numerical examples

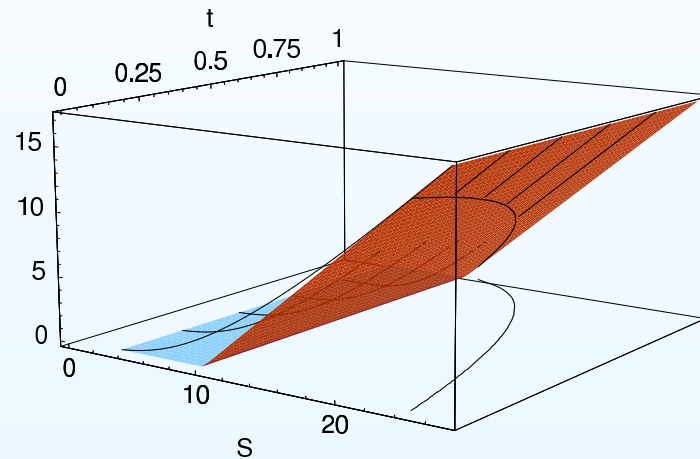
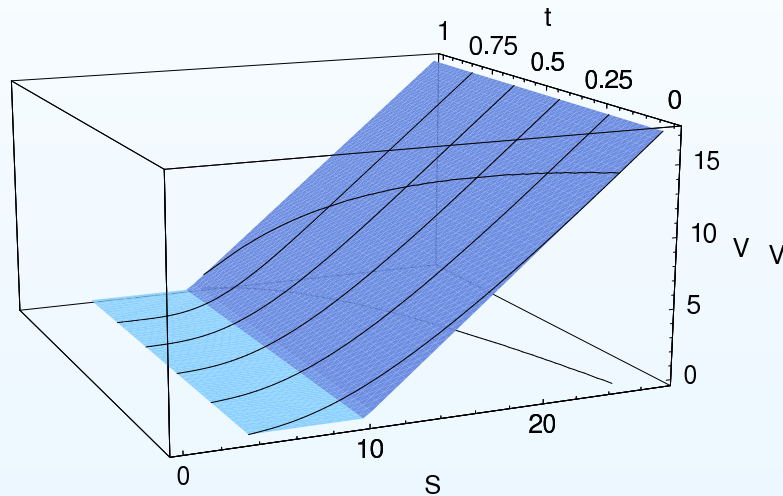
- Pricing American call and put options (for a comparison: price of a European option - dotted line)





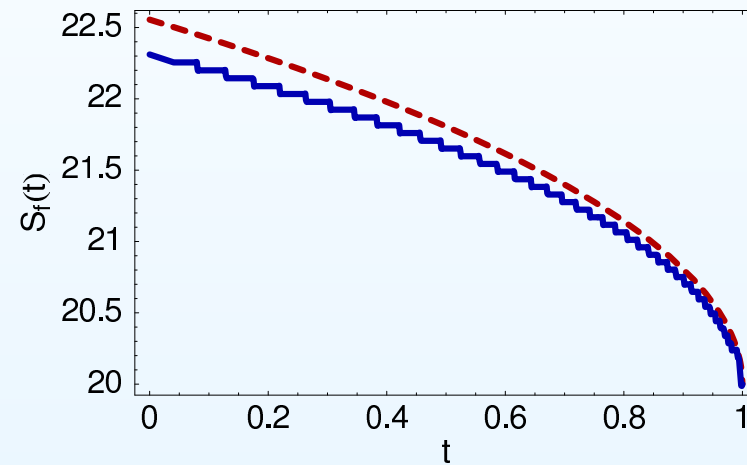
# Numerical examples

- 3D graph for a call option, the free is depicted:



# Numerical examples

- Numerical computation of the free boundary and its comparison with the "square root approximation formula"



# Numerical examples

- M. Lauko, D. Ševčovič: **Comparison of numerical and analytical approximations of the early exercise boundary of American put options**, ANZIAM journal 51, 2010, 430-448.

Comparison of approximation formulae for put options:

